

# Numerik 1 Übungstunde, Gruppe 1 (Donnerstag)

## Woche 02

Übungstunden werden registriert. Bitte informieren Sie den Assistenten, falls Sie nicht in der Aufzeichnung erscheinen möchten.

Streaming (öffentlich) und Aufzeichnungen (Anmeldung nötig):

[https://seminarlive.mnf.uzh.ch/semlive/index.php?id=event\\_detail&module=mat801&semester=fs23&group=1](https://seminarlive.mnf.uzh.ch/semlive/index.php?id=event_detail&module=mat801&semester=fs23&group=1)

Notizen (diese Datei) und Python scripts:

<https://user.math.uzh.ch/florian/stable/page/2023Mat801/notes.php>

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$$S_G^{k,m} \quad \mathbb{P}_m(I)$$

$G$  a "mesh" on  $I$

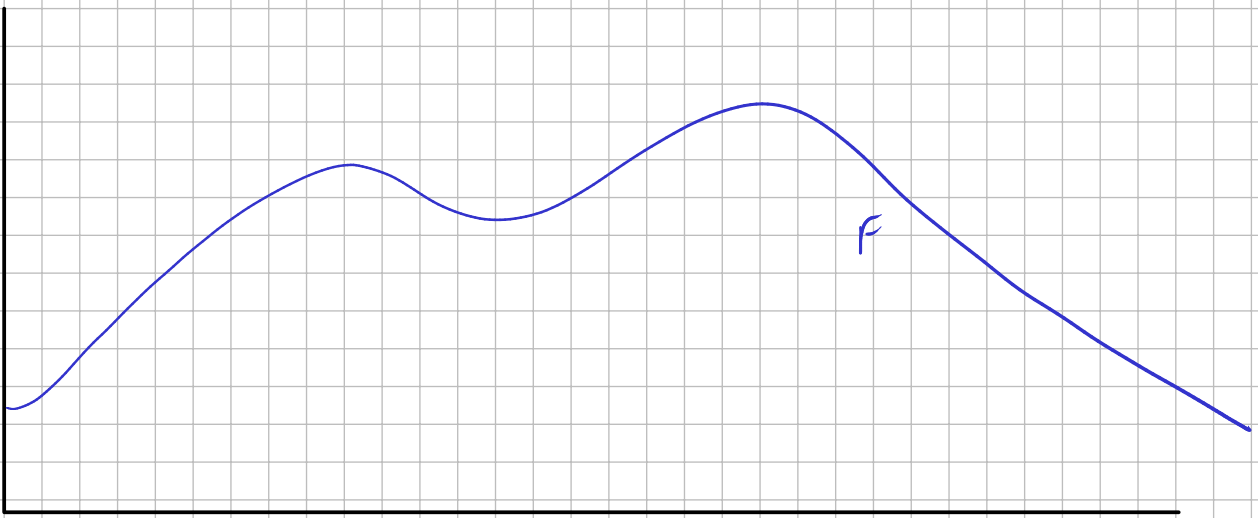
To prove for two sets,  $A, B$

$$A = B$$

You can prove

$$A \subseteq B, \quad B \subseteq A$$

- Take  $f \in S_G^{k,m}$  and prove  $f \in \mathbb{P}_m(I)$ ,
- Take  $h \in \mathbb{P}_m(I)$  " "  $h \in S_G^{k,m}$



We want to interpolate  $F$  "in some way" with "sort of" polynomials such that when increasing the number of nodes, the interpolant gets arbitrarily close to  $F$

- It is clear that not always increasing the degree of the polynomial is enough
- So we use another strategy
  - decrease the interval length
  - Therefore we need to use more than one interval to cover the original domain
  - ( $\Rightarrow$  Splines)

Horner's algorithm

$(x^k)_{k=0}^N$  monomial basis

$$p(x) = \sum_{k=0}^N c_k \cdot b_k(x) = \sum_{k=0}^N c_k x^k$$

## Algorithm 0

Input:  $c$  (coefficients)

$N$

$x$  point where you want to evaluate  $p$

Output:  $p(x)$

evaluation  $= 0$  # this holds  $p(x)$

loop on  $k$  ( $k=0$  to  $N$ )

- compute  $x^k$   $\longrightarrow$
- evaluation  $+ c_k \cdot x^k$

return evaluation

Problem: at step  $k$  we compute  $x^k$   
so at step  $k+1$  we could reuse it \*  
to compute  $x^{k+1}$  as  $\underbrace{x^k \cdot x}_{\text{already computed}}$

## Algorithm 1

Input:  $c$  (coefficients)

$N$

$x$  point where you want to evaluate  $p$

Output:  $p(x)$

evaluation  $= c_0$

polynomial  $= 1$

loop on  $k$ : ( $1$  to  $N$ )

polynomial  $\leftarrow x$

# was  $x^{k-1} \rightarrow x^k = x^{k-1} \cdot x$

evaluation  $+ c_k \cdot \text{polynomial}$

return evaluation

\* We want to compute  $x^3$  (or  $x^{k^m}$ )

1 (start)

$$1 \cdot x \rightarrow x$$

$$x \cdot x \rightarrow x^2$$

$$x^2 \cdot x \rightarrow x^3$$

Algorithm 2 (Horner)

$$\left( \begin{array}{l} c_0 x^0 + c_1 x^1 + c_2 x^2 + c_3 x^3 \\ c_0 + x \cdot (c_1 + x \cdot (c_2 + x \cdot c_3)) \end{array} \right)$$

Input:  $c$  (coefficients)

$N$

$x$  point where you want to evaluate  $p$

Output:  $p(x)$

$$\text{evaluation} = c_N$$

loop on  $n$  (backwards:  $N-1$  to  $0$ ):

$$\text{evaluation} \cdot x = x$$

$$\text{evaluation} + c_n$$

return evaluation

Complexity:

	Steps	complexity per step	total
Algorithm 0:	$O(N)$	$O(N)$	$O(N^2)$
1	$O(N)$	$O(1)$	$O(N)$

$\left( \begin{array}{l} N \\ \sum_{n=0}^N n \end{array} \right)$

2

 $O(N)$  $O(1)$   
(2) $O(N)$ 

Algorithm 2b

Input:  $N$  $x$  (point where you want to evaluate  $p$ ) $s$  (nodes) $b$  (divided differences)Output:  $p(x)$ 

$$p(x) = \sum_{k=0}^N b_k w_{k-1}(x)$$

$$= \sum_{k=0}^N b_k \prod_{l=0}^{k-1} (x - s_l)$$

(monomial basis:  $\prod_{l=0}^{k-1} x$ )evaluation  $n = b_N$ loop over  $k$  ( $N-1 \rightarrow 0$ )evaluation  $+$  =  $x - s_k$ evaluation  $+$  =  $b_k$ 

return evaluation

$$p(x) = b_0 \cdot 1 + b_1 (x - s_0) + b_2 (x - s_0)(x - s_1)$$

$$b_0 + (x - s_0) (b_1 + (x - s_1) \cdot b_2)$$