

Ex 6. Let  $x_1, \dots, x_m \in \mathbb{K}$ . We define for  $k=1, \dots, m$

$$P_k = x_1^k + \dots + x_m^k$$

the  $k$ -th Newton sum.

1. Write a pseudocode that takes as input  $x \in \mathbb{K}$  and an integer  $m$  and returns the list of  $x^k$  for  $k$  power of 2, less than or equal to  $m$  in  $O(\log m)$ .

Pseudocode: **Power**( $x, m$ )

Input:  $x \in \mathbb{K}, m \in \mathbb{Z}$

Output:  $x^0, x^1, x^2, x^4, \dots, x^{2^t}$ , where  $2^t \leq m$

$$t = \lfloor \log_2 m \rfloor$$

$$L[0] = 1, L[1] = x$$

for  $i = 2, \dots, t$  do

$$x = x^2$$

$$L[i] = x$$

end for

The cost is  $O(\log m)$  arithmetic operations since the for is iterated  $\log m$  times.

2. Given  $x_1, \dots, x_m$  returns  $P_k$  in  $O(m \log m)$  arithmetic operations in  $\mathbb{K}$ .

**NewtonSum** ( $m, x_1, \dots, x_m$ )

Input:  $m, L = [x_1, \dots, x_m]$

Output:  $x_1^k + \dots + x_m^k$ , for  $k \leq m$ ,  $k$  power of 2

Sums =  $[0, \dots, 0]$ ; Powers =  $[ ]$ ;

for  $i$  in  $[1, \dots, L]$  do

    Powers[i] = **Powers** ( $x, m$ )

end for

for  $j = 1, \dots, \log_2 m$  do

    for  $k = 1, \dots, m$  do

        Sums[j] += Powers[k][j]

    end for

end for

This will cost  $O(m \log m)$  operations

3. Prove that  $\frac{x}{1-x} = x + x^2 + x^3 + \dots$

If this is true then  $\frac{1}{1-x} = 1 + x + x^2 + \dots$

Call  $\frac{1}{1-x} = S$ .

Then  $S - Sx = 1$  hence

$$S = 1 + Sx$$

$$= 1 + (1 + Sx)x = 1 + x + Sx^2$$

$$= 1 + x + (1 + Sx)x^2 = 1 + x + x^2 + Sx^3$$

$$= \dots \text{ so on } \Rightarrow S = 1 + x + x^2 + \dots$$

4. Deduce that 
$$\sum_{i=1}^m \frac{x_i X}{1 - x_i X} = \sum_{k \geq 1} P_k X^k$$

To avoid confusion, instead of  $x_i$  write  $\alpha_i$

$$\frac{\alpha_i X}{1 - \alpha_i X} = (\alpha_i X) + (\alpha_i X)^2 + (\alpha_i X)^3 + \dots$$

$$\text{Hence } \sum_{i=1}^m \frac{\alpha_i X}{1 - \alpha_i X} = \sum_{i=1}^m (\alpha_i X + (\alpha_i X)^2 + \dots)$$

$$= \sum_{i=1}^m \left( \sum_{k \geq 1} (\alpha_i X)^k \right) =$$

$$= \sum_{k \geq 1} \left( \sum_{i=1}^m (\alpha_i)^k \right) X^k = \sum_{k \geq 1} P_k X^k$$

5. Compute  $P_k$  for  $k \leq m$  in  $O(M(m) \log m)$

$$\text{Since } \sum_{k \geq 1} P_k X^k = \sum_{i=1}^m \frac{\alpha_i X}{1 - \alpha_i X}$$

we compute the sums of  $\frac{A_i(x)}{B_i(x)}$  where

$$A_i(x) = \alpha_i x \quad \text{and} \quad B_i(x) = 1 - \alpha_i(x)$$

with a divide and conquer approach,